

# Recognisable sets, profinite topologies and weak arithmetic

Jean-Éric Pin

IRIF, CNRS and University Paris 7



May 2018, Firenze

Partially funded by the ERC (grant agreement No 670624) and  
by the DeLTA project (ANR-16-CE40-0007) 

# Outline

- (1) Original motivation
- (2) Recognisable sets
- (3) Topological characterizations
- (4) Functions from  $\mathbb{N}$  to  $\mathbb{N}$
- (5) Transductions

# The starting point

If  $A$  is a **one-letter** alphabet, the **free monoid**  $A^*$  is isomorphic to the additive monoid  $\mathbb{N}$ .

It seems natural to extend results on  $\mathbb{N}$  to  $A^*$ . However, one may expect any result on  $A^*$  to become **trivial** on a **one-letter** alphabet.

Surprisingly enough, this is **not always** the case...

# An example

Given a language  $L \subseteq A^*$  and a word  $u \in A^*$ , let

$$u^{-1}L = \{x \in A^* \mid ux \in L\}$$

$$Lu^{-1} = \{x \in A^* \mid xu \in L\}$$

**Theorem** (Almeida, Esik, Pin 2017)

*A class of regular languages closed under finite intersection, finite union, quotients and inverse of length-decreasing morphisms is also closed under inverse of morphisms.*

# For one-letter alphabets

For  $L \subseteq \mathbb{N}$  and  $k > 0$ , let

$$L - 1 = \{n \in \mathbb{N} \mid n + 1 \in L\}$$

$$L \div k = \{n \in \mathbb{N} \mid kn \in L\}$$

**Corollary** (Cegielski, Grigorieff, Guessarian 2014)

Let  $\mathcal{L}$  be a *lattice of regular subsets* of  $\mathbb{N}$  such that if  $L \in \mathcal{L}$ , then  $L - 1 \in \mathcal{L}$ . Then for each positive integer  $k$ ,  $L \in \mathcal{L}$  implies  $L \div k \in \mathcal{L}$ .

## Corollary

The mapping

$$\mathcal{V} \mapsto \{X_L : L \in \mathcal{V}(a^*)\}$$

is an isomorphism from the lattice of commutative positive (*ld*-)varieties to the sublattice of  $\mathcal{P}(\mathcal{P}(\mathbb{N}))$  consisting of those sets  $\mathcal{X}$  of finite or ultimately periodic subsets of  $\mathbb{N}$  that contain  $\emptyset$  and  $\mathbb{N}$  which are closed under union and intersection, moreover, the *decrement operation* defined by

$$X \rightarrow X - 1 = \{n - 1 \mid n \in X, n > 0\}.$$

Any such set  $\mathcal{X}$  is closed under the *division operations* defined by:

$$X \rightarrow X/d = \{n \mid nd \in X\}, \quad d > 1.$$

When restricted to commutative (*ld*-)varieties, the same mapping creates an order isomorphism from the lattice of commutative (*ld*-)varieties to the sublattice of  $\mathcal{P}(\mathcal{P}(\mathbb{N}))$  consisting of all those sets  $\mathcal{X}$  of finite or ultimately periodic subsets of  $\mathbb{N}$  which are additionally closed under complementation.

# Original motivation

A function  $f : A^* \rightarrow B^*$  is **regularity-preserving** if, for each regular language  $L$  of  $B^*$ ,  $f^{-1}(L)$  is also **regular**.

More generally, let  $\mathcal{C}$  be a class of **regular languages**. A function  $f : A^* \rightarrow B^*$  is  **$\mathcal{C}$ -preserving** if, for each  $L \in \mathcal{C}$ ,  $f^{-1}(L)$  is also in  $\mathcal{C}$ .

**Goal.** Find a complete description of **regularity-preserving** [ **$\mathcal{C}$ -preserving**] functions.

Same questions for **transductions**, that is, relations from  $A^*$  to  $B^*$ .

# Part I

## Recognisable sets

# Monoids

A **monoid** is a set  $M$  equipped with an associative binary operation (the **product**) and an identity  $1$  for this operation.

A monoid  $M$  is **finitely generated** if there exists a finite subset  $F$  of  $M$  which generates  $M$ .

**Examples.** The **free monoid**  $A^*$ , with  $A$  finite.

Given a monoid  $M$ , the set  $\mathcal{P}(M)$  of **subsets** of  $M$  is a monoid under the product defined, for  $X, Y \subseteq M$ , by  $XY = \{xy \mid x \in X, y \in Y\}$ .

# Recognisable subsets of a monoid

A subset  $P$  of a monoid  $M$  is **recognizable** if there exists a **finite monoid**  $F$ , a monoid morphism  $\varphi : M \rightarrow F$  and a subset  $Q$  of  $F$  such that  $P = \varphi^{-1}(Q)$ .

$\text{Rec}(M)$  = set of **recognizable subsets** of  $M$  .

## Theorem (Kleene)

*If  $M = A^*$ , then **recognizable** = **rational** = **regular** (that is, recognised by a finite automaton).*

# Recognisable subsets of $\mathbb{N}$

An **arithmetic progression** is a subset of  $\mathbb{N}$  of the form  $a + r\mathbb{N}$ , with  $r \geq 0$ .

A subset of  $\mathbb{N}$  is **recognizable iff** it is a **finite union of arithmetic progressions**.

$\{1, 3, 4, 7, 8, 9, 11, 12, 13, 17, 18, 22, 23, 27, 28, \dots\} =$   
 $\{1, 3, 4, 9, 11\} \cup \{7 + 5n \mid n \geq 0\} \cup \{8 + 5n \mid n \geq 0\}$   
is a finite union of arithmetic progressions.

# Recognisable subsets of a product of monoids

## Theorem (Mezei)

Let  $M = M_1 \times \cdots \times M_n$  be a product of monoids. A subset of  $M$  is *recognisable* iff it is a *finite union* of subsets of the form  $R_1 \times \cdots \times R_n$ , where each  $R_i$  is a *recognisable* subset of  $M_i$ .

**Exercise:** find the *recognisable* subsets of  $\mathbb{N}^k$ .

# Transductions

Given two monoids  $M$  and  $N$ , a **transduction** from  $M$  into  $N$  is a relation on  $M$  and  $N$ .

If  $\tau : M \rightarrow N$  is a **transduction**, then the **inverse relation**  $\tau^{-1} : N \rightarrow M$  is also a **transduction**. If  $R \subseteq N$ , then

$$\tau^{-1}(R) = \{x \in M \mid \tau(x) \cap R \neq \emptyset\}$$

A function  $f : M \rightarrow N$  is **recognizability-preserving** if, for each  $R \in \text{Rec}(N)$ ,  $f^{-1}(R) \in \text{Rec}(M)$ .

Similarly,  $\tau : M \rightarrow N$  is **recognizability-preserving** if, for each  $R \in \text{Rec}(N)$ ,  $\tau^{-1}(R) \in \text{Rec}(M)$ .

# Part II

## Topological characterizations

# Residually finite monoids

Let  $M$  be a monoid. A monoid  $F$  separates two elements  $x, y \in M$  if there exists a morphism  $\varphi : M \rightarrow F$  such that  $\varphi(x) \neq \varphi(y)$ .

A monoid is **residually finite** if any pair of distinct elements of  $M$  can be separated by a **finite monoid**.

Finite monoids, free monoids, free groups are **residually finite**. A product of **residually finite** monoids is **residually finite**.

# Profinite metric

Let  $M$  be a residually finite monoid. The **profinite metric**  $d$  is defined by setting, for  $u, v \in M$ :

$$r(u, v) = \min\{|F| \mid F \text{ is a monoid separating } u \text{ and } v\}$$

$$d(u, v) = 2^{-r(u, v)}$$

with  $\min \emptyset = +\infty$  and  $2^{-\infty} = 0$ . Then

$$d(u, w) \leq \max(d(u, v), d(v, w)) \quad (\text{ultrametric})$$

$$d(uw, vw) \leq d(u, v)$$

$$d(wu, wv) \leq d(u, v)$$

# Recognizability-preserving functions

Let  $M$  and  $N$  be two finitely generated, residually finite monoids. (For instance  $M = A^*$  and  $N = B^*$ ).

## Theorem (Pin-Silva 2005)

A function  $M \rightarrow N$  is *recognizability-preserving* iff it is *uniformly continuous*.

## Proposition (Pin-Silva 2005)

The function  $\tau : M \times \mathbb{N} \rightarrow M$  defined by  $\tau(x, n) = x^n$  is *recognizability-preserving*.

**Corollary.** The function  $u \rightarrow u^{|u|}$  (from  $A^*$  to  $A^*$ ) is *recognizability-preserving*. Indeed it can be decomposed as

$$A^* \rightarrow A^* \times \mathbb{N}$$

$$u \rightarrow (u, |u|)$$

$$A^* \times \mathbb{N} \rightarrow A^*$$

$$(u, n) \rightarrow u^n$$

# Some examples of regularity preserving functions

$$u \rightarrow u^2$$

$$u \rightarrow \tilde{u}u$$

$$u \rightarrow u^{|u|}$$

$$u \rightarrow a^{|u|_a} b^{|u|_b}$$

$$a^m c b^n \rightarrow a^n b^{mn}$$

$$u_0 \# u_1 \# u_2 \rightarrow u_2 \# u_1 \# u_0 \# u_1 \# u_2$$

# Part III

## Functions from $\mathbb{N}$ to $\mathbb{N}$

# Ultimately periodic functions

A function  $f : \mathbb{N} \rightarrow \mathbb{N}$  is **ultimately periodic** if there exists  $t \geq 0$  and  $p > 0$  such that, for all  $n \geq t$ ,  
 $f(n + p) = f(n)$ .

A function  $f : \mathbb{N} \rightarrow \mathbb{N}$  is **ultimately periodic modulo  $n$**  if the function  $f \bmod n$  is ultimately periodic.

A function  $f : \mathbb{N} \rightarrow \mathbb{N}$  is **cyclically ultimately periodic** if it is ultimately periodic modulo  $n$  for all  $n > 0$ .

# Regularity-preserving functions from $\mathbb{N}$ to $\mathbb{N}$

**Theorem** (Siefkes 1970, SeiferasMcNaughton 1976)

A function  $f : \mathbb{N} \rightarrow \mathbb{N}$  is *ultimately periodic modulo  $n$*  iff for  $0 \leq k < n$ , the set  $f^{-1}(k + n\mathbb{N})$  is *regular*.

**Theorem** (Siefkes 1970, SeiferasMcNaughton 1976)

A function  $f : \mathbb{N} \rightarrow \mathbb{N}$  is *regularity-preserving* iff it is *cyclically ultimately periodic* and, for every  $k \in \mathbb{N}$ , the set  $f^{-1}(k)$  is *regular*.

# Ultimately periodic functions

A function  $f : \mathbb{N} \rightarrow \mathbb{N}$  is **ultimately periodic modulo  $k$**  if the function  $f \bmod k$  is **ultimately periodic**.

It is **cyclically ultimately periodic (cup)** if it is ultimately periodic modulo  $n$  for all  $n > 0$ .

**Proposition** (Siefkes 70, SeiferasMcNaughton 76)

*A function  $f : \mathbb{N} \rightarrow \mathbb{N}$  is **ultimately periodic modulo  $n$**  iff for  $0 \leq k < n$ , the set  $f^{-1}(k + n\mathbb{N})$  is regular. It is **regularity-preserving** iff it is **cyclically ultimately periodic** and  $f^{-1}(k)$  is **regular** for every  $k \in \mathbb{N}$ .*

# Two examples

## Theorem (Siefkes 1970)

The functions  $n \rightarrow 2^n$  and  $n \rightarrow 2^{2^{2^{\dots^2}}}$  (exponential stack of 2's of height  $n$ ) are *cyclically ultimately periodic* and hence *regularity-preserving*.

## Theorem (Siefkes 70, Zhang 98, Carton-Thomas 02)

Let  $f, g : \mathbb{N} \rightarrow \mathbb{N}$  be *cyclically ultimately periodic functions*. Then so are the following functions:

- (1)  $g \circ f$ ,  $f + g$ ,  $fg$ ,  $f^g$ , and  $f - g$  provided that  $f \geq g$  and  $\lim_{n \rightarrow \infty} (f - g)(n) = +\infty$ ,
- (2) (*generalised sum*)  $n \rightarrow \sum_{0 \leq i \leq g(n)} f(i)$ ,
- (3) (*generalised product*)  $n \rightarrow \prod_{0 \leq i \leq g(n)} f(i)$ .

## Two counterexamples

[Siefkes 1970] The function  $n \rightarrow \lfloor \sqrt{n} \rfloor$  is **not** cyclically ultimately periodic and hence **not** regularity-preserving.

The function  $n \rightarrow \binom{2n}{n}$  is **not** ultimately periodic modulo 4 and hence **not** regularity-preserving. Indeed

$$\binom{2n}{n} \bmod 4 = \begin{cases} 2 & \text{if } n \text{ is a power of } 2, \\ 0 & \text{otherwise.} \end{cases}$$

# Recursivity

Let  $f : \mathbb{N} \rightarrow \{0, 1\}$  be a **non-recursive** function.  
Then the function  $n \rightarrow (\sum_{0 \leq i \leq n} f(i))!$  is  
**regularity-preserving** but **non-recursive**.

**Open problem.** Is it possible to describe all  
**recursive regularity-preserving** functions, respectively  
all **recursive cyclically ultimately periodic** functions?

One could try to use **Siefkes' recursion scheme**  
(1970).

## Theorem

Let  $g : \mathbb{N}^k \rightarrow \mathbb{N}$  and  $h : \mathbb{N}^{k+2} \rightarrow \mathbb{N}$  be *cyclically ultimately periodic* functions satisfying three technical conditions. Then the function  $f$  defined from  $g$  and  $h$  by *primitive recursion*, i.e.

$$f(0, x_1, \dots, x_k) = g(x_1, \dots, x_k),$$

$$f(n + 1, x_1, \dots, x_k) = h(n, x_1, \dots, x_k, f(n, x_1, \dots, x_k))$$

is *cyclically ultimately periodic*.

# The three technical conditions

- (1)  $h$  is cyclically ultimately periodic in  $x_{k+2}$  of decreasing period,
- (2)  $g$  is essentially increasing in  $x_k$ ,
- (3) for all  $x \in \mathbb{N}^{k+2}$ ,  $x_{k+2} < h(x_1, \dots, x_{k+2})$ .

A function  $f$  is essentially increasing in  $x_j$  iff, for all  $z \in \mathbb{N}$ , there exists  $y \in \mathbb{N}$  such that for all  $x \in \mathbb{N}^n$ ,  $y \leq x_j$  implies  $z \leq f(x_1, \dots, x_n)$ .

A function  $f$  is c.u.p. of decreasing period in  $x_j$  iff, for all  $p$ , the period of the function  $f \bmod p$  in  $x_j$  is  $\leq p$ .

# Part IV

## An extension

# Lattice of subsets

Let  $X$  be a set. A **lattice of subsets of  $X$**  is a set  $\mathcal{L}$  of subsets of  $X$  containing  $\emptyset$  and  $X$  and closed under finite **union** and finite **intersection**.

A **Boolean algebra of subsets of  $X$**  is a lattice of subsets of  $X$  closed under complement.



A **Pervin space** is a pair  $(X, \mathcal{L})$  where  $\mathcal{L}$  is a lattice of subsets of  $X$ .

# Lattice-preserving functions

Let  $f : X \rightarrow Y$  be a map,  $\mathcal{K}$  be a lattice of subsets of  $X$  and  $\mathcal{L}$  a lattice of subsets of  $Y$ .

## Theorem

*The following conditions are equivalent:*

- (1) *for each  $L \in \mathcal{L}$ ,  $f^{-1}(L) \in \mathcal{K}$ ,*
- (2)  *$f$  is a uniformly continuous map from  $(X, \mathcal{K})$  to  $(Y, \mathcal{L})$ .*

# Lattice-preserving functions

Let  $f : X \rightarrow Y$  be a map,  $\mathcal{K}$  be a lattice of subsets of  $X$  and  $\mathcal{L}$  a lattice of subsets of  $Y$ .

## Theorem

*The following conditions are equivalent:*

- (1) *for each  $L \in \mathcal{L}$ ,  $f^{-1}(L) \in \mathcal{K}$ ,*
- (2)  *$f$  is a uniformly continuous map from  $(X, \mathcal{K})$  to  $(Y, \mathcal{L})$ .*

Wait a second, what does **uniformly continuous** mean in this setting?

# Uniform spaces

A **uniformity** on a set  $X$  is a nonempty set  $\mathcal{U}$  of **reflexive** relations (**entourages**) on  $X$  such that:

- (1) if a relation  $U$  on  $X$  contains an element of  $\mathcal{U}$ , then  $U \in \mathcal{U}$ , (**extension property**),
- (2) the intersection of any two elements of  $\mathcal{U}$  is in  $\mathcal{U}$ , (**intersection**),
- (3) for each  $U \in \mathcal{U}$ , there exists  $V \in \mathcal{U}$  such that  $VV \subseteq U$  (**sort of transitivity**).
- (4) for each  $U \in \mathcal{U}$ ,  ${}^tU \in \mathcal{U}$  (**symmetry**).

# Quasi-uniform spaces

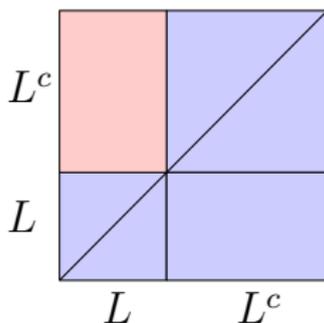
A **quasi-uniformity** on a set  $X$  is a nonempty set  $\mathcal{U}$  of **reflexive** relations (**entourages**) on  $X$  such that:

- (1) if a relation  $U$  on  $X$  contains an element of  $\mathcal{U}$ , then  $U \in \mathcal{U}$  (**extension property**),
- (2) the intersection of any two elements of  $\mathcal{U}$  is in  $\mathcal{U}$  (**intersection**),
- (3) for each  $U \in \mathcal{U}$ , there exists  $V \in \mathcal{U}$  such that  $VV \subseteq U$  (**sort of transitivity**).

# Pervin spaces as quasi-uniform spaces

Let  $(X, \mathcal{L})$  be a Pervin space. For each  $L \in \mathcal{L}$ , let

$$\begin{aligned} U_L &= (X \times L) \cup (L^c \times X) \\ &= \{(x, y) \in X \times X \mid x \in L \Rightarrow y \in L\} \end{aligned}$$



The sets  $U_L$  form the subbasis of a **quasi-uniformity**.

# Uniform continuity

Let  $X$  and  $Y$  be quasi-uniform spaces. A function  $f: X \rightarrow Y$  is uniformly continuous if, for each entourage  $V$  of  $Y$ ,  $(f \times f)^{-1}(V)$  is an entourage of  $X$ .

## Proposition

Let  $(X, \mathcal{K})$  and  $(Y, \mathcal{L})$  be two Pervin spaces. A function  $f: X \rightarrow Y$  is uniformly continuous iff for each  $L \in \mathcal{L}$ ,  $f^{-1}(L) \in \mathcal{K}$ .

# Generalized ultrametric

A **generalized ultrametric** on a set  $X$  is a mapping  $d : X \times X \rightarrow \mathbb{R}^+$  satisfying the following conditions:

- (1) for all  $x \in X$ ,  $d(x, x) = 0$ .
- (2) for all  $x, y, z \in X$ ,  
 $d(x, z) \leq \max(d(x, y), d(y, z))$ .

Let  $(X, \mathcal{L})$  be a Pervin space. Are equivalent:

- (1) The associated **quasi-uniformity** can be defined by a generalized ultrametric,
- (2) The **quasi-uniformity** has a **countable basis**,
- (3) The lattice  $\mathcal{L}$  is **countable**.

# Boolean algebras

If  $\mathcal{L}$  is a Boolean algebras, then one has a **uniformity**. Moreover if  $\mathcal{L}$  is **countable**, this uniformity can be defined by an **ultrametric**.

If  $\mathcal{L}$  is the set of **recognizable subsets** of a residually finite monoid  $M$ , then this ultrametric is the **profinite ultrametric**.

# Part V

## Transductions

# Recognizability-preserving transductions

Let  $M$  and  $N$  be two finitely generated, residually finite monoids.

## Theorem

A function  $M \rightarrow N$  is *recognizability-preserving* iff it is *uniformly continuous*.

What about *transductions* from  $M$  to  $N$ ?

# Completion

Let  $M$  be a **finitely generated, residually finite** monoid. Let  $\widehat{M}$  be the **completion** of the metric space  $(M, d)$ .

## Proposition

$\widehat{M}$  is a **compact monoid**.

Moreover, the set  $\mathcal{K}(\widehat{M})$  of **compact** subsets of  $\widehat{M}$  is also a compact monoid for the **Hausdorff metric**.

# Back to transductions

Let  $M$  and  $N$  be two finitely generated, residually finite monoids and let  $\tau : M \rightarrow N$  be a transduction.

Define a map  $\hat{\tau} : M \rightarrow \mathcal{K}(\hat{N})$  by setting, for each  $x \in M$ ,  $\hat{\tau}(x) = \overline{\tau(x)}$ .

## Theorem (Pin-Silva 2005)

*The transduction  $\tau$  is recognizability-preserving iff  $\hat{\tau}$  is uniformly continuous.*

# Exercises

Let  $L$  be a subset of  $A^*$ . Let

$$\frac{1}{2n+1}L = \{u \in A^* \mid \text{there exist } x, y \in A^*, \\ |x| = |y| = n \text{ and } xuy \in L\}$$

If  $L$  is regular, then so is the language

$$\bigcup_{p \text{ odd prime}} \frac{1}{p}L$$

The transduction  $u \rightarrow u^*$  is regularity-preserving.

# Part VI

## $p$ -group languages

**Target class  $\mathcal{G}_p$ :** the class of languages recognized by a finite  $p$ -group.

**Goal.** Characterization of  $\mathcal{G}_p$ -preserving functions.

# Functions from $\mathbb{N}$ to $\mathbb{Z}$

The **difference operator**  $\Delta$  associates to each function  $f : \mathbb{N} \rightarrow \mathbb{Z}$ , the function  $\Delta f : \mathbb{N} \rightarrow \mathbb{Z}$  defined by  $(\Delta f)(n) = f(n+1) - f(n)$ .

A **Newton polynomial** is a function  $f$  such that  $\Delta^k f = 0$  for almost all  $k$ .

# Mahler's theorem

Let  $\delta^k f = (\Delta^k f)(0)$ .

## Theorem (Mahler 58)

Let  $f : \mathbb{N} \rightarrow \mathbb{Z}$  be a function. Are equivalent:

- (1)  $f$  is *uniformly continuous* for the  $p$ -adic metric,
- (2) the functions  $\Delta^n f$  tend uniformly to 0,
- (3) the  $p$ -adic norm of  $\delta^n f$  tends to 0,
- (4)  $f$  is the *uniform limit* of a sequence of Newton polynomials.